

Overflowing Attack Potential

Scoring Øefence-in-Depth

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09 10 af 0f a1 b2 da 29 89 a8 82 8a 9b bc 90 c3 87 90 90 a4 9f ff a4 89 8a 53 46 89













- **1.Buffer overflows, a bit of background**
- 2.Reviewing and bypassing defence-indepth techniques
- **3.Impact in the CC**
- 4.What to do?

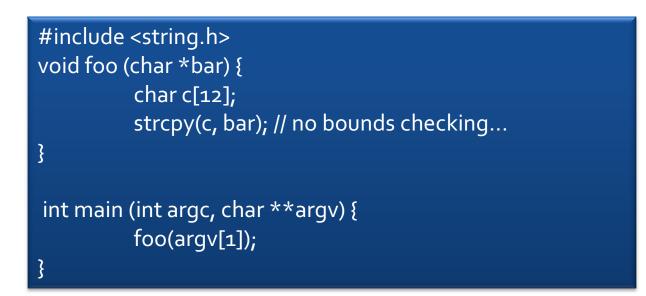


2.Reviewing and bypassing defence-indepth techniques

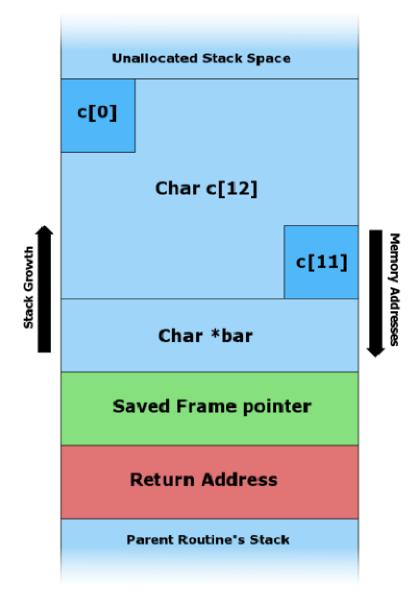
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- 4.What to do?



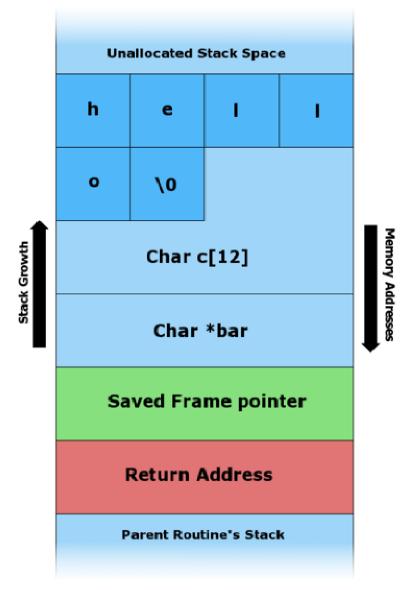
You know... The classic stack overflow....



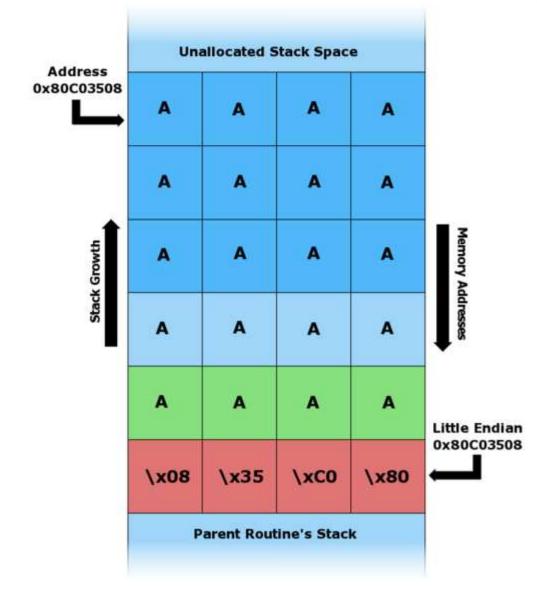














There are also heap and integer overflows....



Could lead to arbitrary code execution





Those were the old days...

Very few problems for the attacker:

- Null bytes
- Shellcode size and other constraints
- Shellcode development



- **3.Impact in the CC**
- 4.What to do?



- Stack canaries approach:
 - The compiler place a value before the return address when a function is called and check that the value has not changed when the function finalize.





- Bypassing stack canaries:
 - Implementation can be not correct
 - It can be a statistical problem



Bypassing stack canaries:

Windows: SEH overwriting Protected by SafeSEH and so on... Unix: Other (more complex)

techniques...



- Non-eXecutable Stack approach:
 - Effective implementation in hardware. Widely deployed (every computer since 2001 allow this)
 - Code is code and data is data
 - However, it is easy to bypass



- Bypassing Non-eXecutable Stack:
 - Save the payload in the heap
 - Return into libc (standard C library) attacks

64bits hardware saves the way:



EPOCHE&ESPRI

- In 64 bits personal computers, arguments are loaded in registers, not in the stack.
- Return into libc attack is not possible



64bits hardware saves the way:



Return oriented programming



Return on Enter Programming SLIKE AT SOM TOTE, BUT instead Of Cutting Cut Lewers magazines you are Culting Out INSTRUCTIONS FOM EX SEGMENS



 ASLR (Address Space Layout Randomization) approach:

 The code is loaded in different memory regions each time

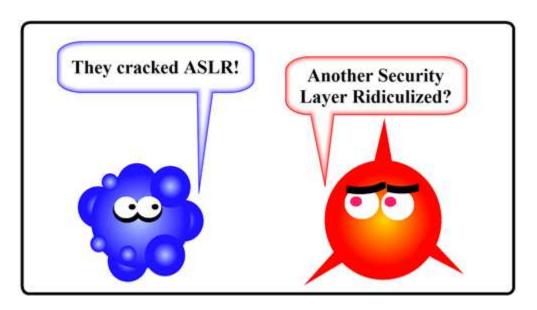


- Bypassing ASLR:
 - It could be an statistical question





Bypassing ASLR: Maybe not all the libraries are randomly loaded





- Mixed approach:
 - Standalone use of this techniques is not very useful

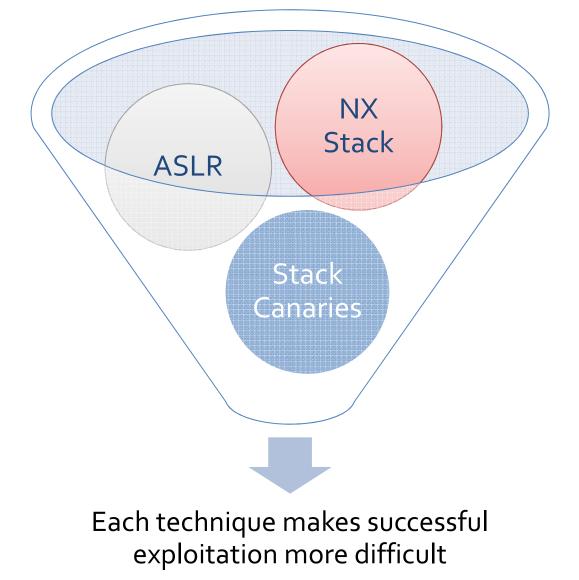




Non-eXecutable Stack + ASLR:

Make very difficult the return attacks.

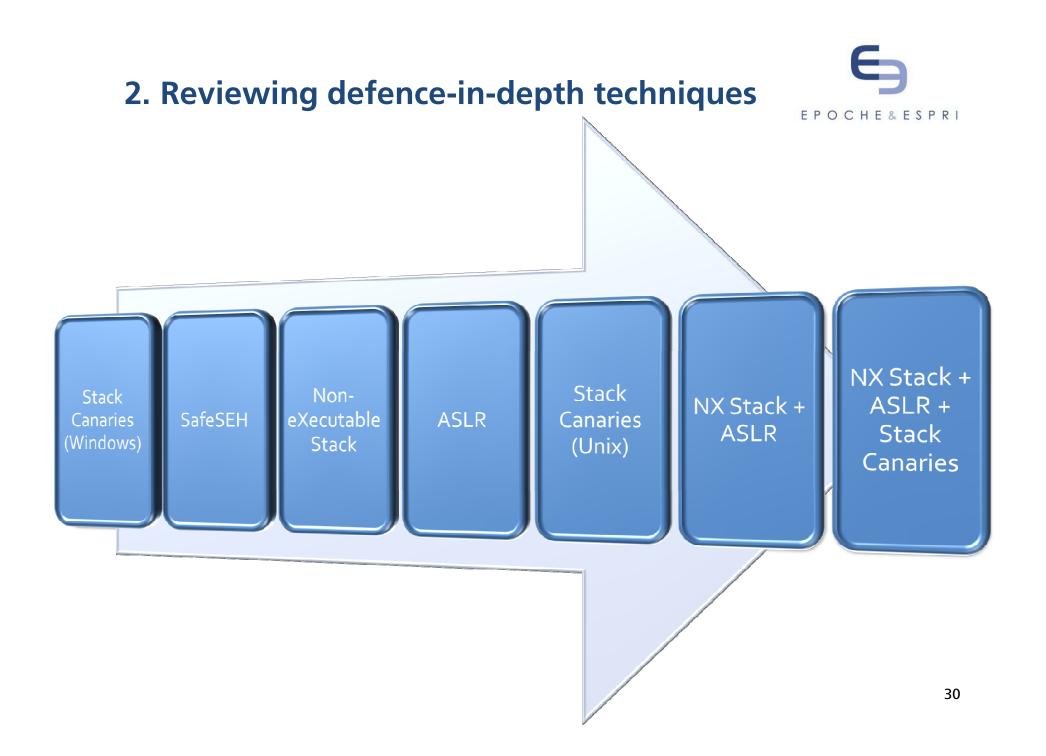






There exists more defence-in-depth techniques

Attackers also develop new techniques to bypass the countermeasures









2.Reviewing defence-in-depth techniques

3.Impact in the CC

4.What to do?

3. Impact in the CC

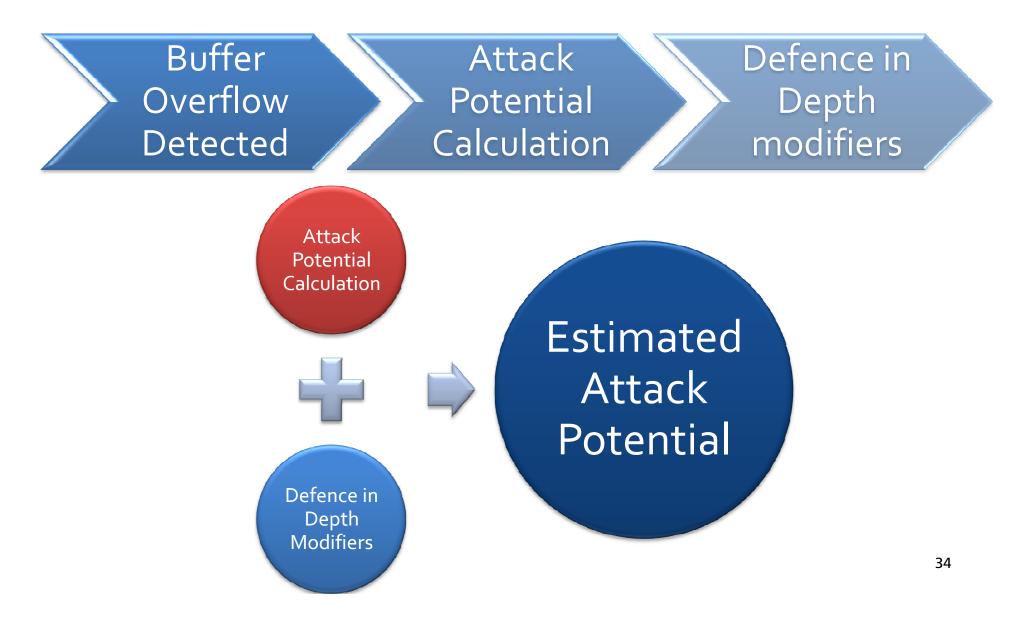


- We start from a detected buffer overflow
 - Unique characteristics
 - Unique exploit path

Attack potential calculation

3. Impact in the CC







3. Impact in the CC

Defence in depth technique	Attack potential factor
Stack Canaries (Windows)	X 1.2
SafeSEH	x 1.3
Non-eXecutable Stack	× 1.35
ASLR	X 1.50
Stack Canaries (Unix)	X 1.52
NX Stack + ASLR	× 1.54
NX Stack + ASLR + Stack Canaries (Windows)	x 1.62
NX Stack + ASLR + Stack Canaries (Windows) + SafeSEH	x 1.66
NX Stack + ASLR + Stack Canaries (Linux)	x 1.68

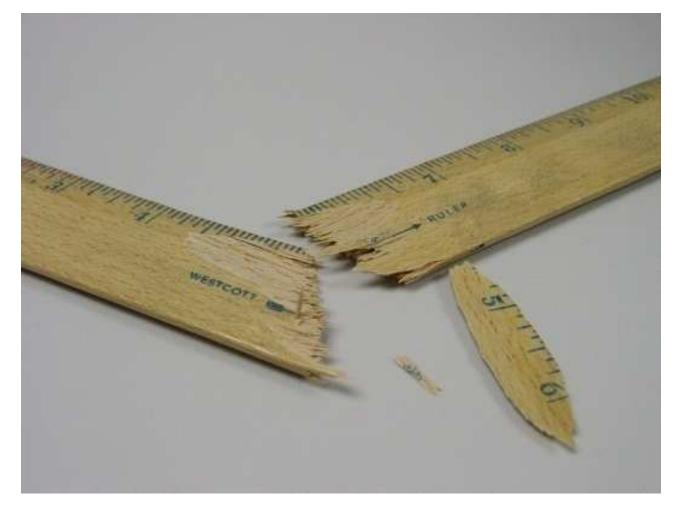


2.Reviewing defence-in-depth techniques

3.Impact in the CC







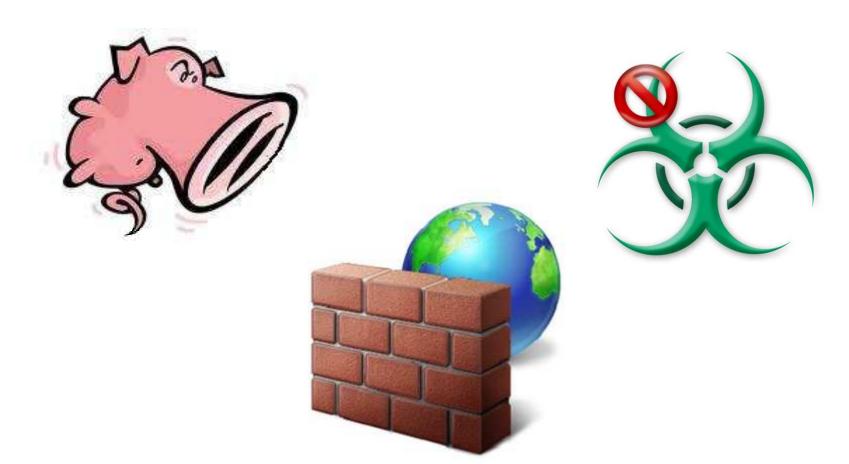
THERE IS NO EXACT RULE





TIME CHANGES THE THINGS





OTHER AREAS OF APPLICATION?



Apply those techniques!

- Whenever it is possible
- Through compiler
- Through Operating System







Thanks for your attention!

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